

DESIGN AND IMPLEMENTATION OF IDENTITY DISTRIBUTED DECRYPTION SCHEME FOR E- HEALTH RECORDS

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Abstract – The fast improvement of the Internet of Things (IoT) has incited the advancement of a steadily expanding number of novel applications recently. One of them is the e-prosperity system, which can outfit people with first class and supportive clinical consideration. Meanwhile, it is a primary concern of conflict and hardships to defend the assurance and security of the client's own personal prosperity record. A couple of cryptographic procedures have been proposed, for instance, scramble client's data preceding sharing it. Anyway, it is tangled to give the data to different social events (subject matter experts, prosperity divisions, etc), due to the way that data should be mixed under each recipient's keys. But a couple (t; n) edge secret sharing plans can share the data simply need one encryption movement, there is a requirement that the disentangling private key should be recreated by one party. To offset this insufficiency, in this paper, we propose a useful person based scattered translating plan for individual prosperity record sharing structure. It is invaluable to give their data to various get-togethers and doesn't require reproducing the interpreting private key.

Index terms – Internet of things, IBE, PHR, Public Key Infrastructure.

I. INTRODUCTION

E-prosperity systems bring clients a lot of benefits. They can save lives in emergency clinical conditions, through the steady checking of the related contraptions; it is easy to perceive the emergency conditions, for instance, asthma attacks, cardiovascular breakdown and diabetes. As shown in Fig 1,

the clinical data and prosperity data are assembled by the related contraptions. Then, the data is moved to the subject matter expert or the clinical consideration office by distant association devices, similar to cells and tablets. Honestly, these data are significant for individual prosperity records (PHRs).



Fig. 1: Data Collect by Smartphone

PHR consolidates prosperity data, yet what's more a few critical information associated with patient thought. This data is managed by the patient and ordinarily set aside in the cloud server (clinical servers) [4]. Unlike the electronic clinical record, the PHR isn't made and stayed aware of by establishments (like clinical establishments and facilities). The data collection and move ing are done by the patient. The justification for PHRs is to store an exact and finish once-over of the solitary's clinical history. They license endorsed clients or establishments to get to the data over the Internet.

Another audit [5] shows that a larger piece of clients use applications and various instruments to follow their wellbeing, food and rest; 44% of people have gotten to their clinical records on the web. Like in Fig 2, in an ordinary e-prosperity individual prosperity records (PHRs) designing, the client's data are assembled and delivered off the clinical servers. At the point when the expert necessities to review the client's PHRs

(clinical data, clinical record, etc), he truly needs to download the PHRs from the clinical server. Incidentally, the immense PHRs data are typically taken care of and dealt with in the cloud stage, for instance, Amazon Web Services, Google Cloud. In view of the PHRs contains some sensitive and high-regard data; the cloud server has transformed into an engaging goal for hacking.

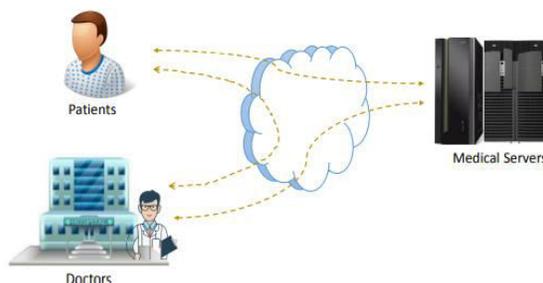


Fig. 2: e-health PHR architecture

Undoubtedly, guaranteeing the security of the patient's PHRs is basic. According to the latest report by Protenus [6], more than 31 million patients of prosperity records have proactively been infiltrated in the key portion of 2019, with hacking causing the majority of security episodes and breaking the most grasping records. As referred to in [7], a few mystery key based designs and electronic mark plans taking into account Public Key Infrastructure (PKI) have been proposed for affirmation. Various public key encryption plans [8], [9], [10] have in like manner been proposed to defend client's PHRs. In case of a data break, the security or insurance of the PHRs data will

not be compromised, due to truth that the adversary can't get the deciphering private key.

II. LITERATURE SURVEY

In the traditional PKI framework, managing the enormous public keys is tangled. The person based cryptography (IBC) offers a cunning choice. Potentially the most charming property is that the public key of a component is its character. Additionally, in the IBC structures, the supports are not commonly needed. Overall, the public key can be figured from its personality string by a predefined computation, (for instance, a hash work) for specific commitments of public limits.

Boneh and Franklin introduced the essential character based encryption plan from pairings. In their arrangement, an accepted party named Key Generation Center (KGC) is involved, which can remove client's private key by using its master secret key and the character of the client. Furthermore, the public key is a hash worth of the character, and can be used to encode the messages. Following Boneh's thought, different IBE plans have been proposed.

Many existed IBE plans rely upon bundles with the bilinear aide of different assumptions. In 2004, Boneh and Boyen introduced a substitute IBE plan, it might be exhibited secure without using the sporadic prophet

model. They similarly proposed two explicit character secure IBE plans, while these plans were furthermore secure under the inconsistent prophet model. An extended variation was presented which is a particular person Hierarchical IBE plot. Waters proposed the fundamental IBE plot, which is totally secure without erratic prophets.

Some IBE plans taking into account the quadratic residuosity assumption have been proposed. In the composing, the size of the ciphertext is short, under the quadratic residuosity issue, the arrangement was exhibited secure. In any case, the encryption and it are not adequately useful to interpret plans. Lately, a couple of plans considering the learning-with-botches (LWE) assumption safe house been proposed. In makers showed a strategy for building stowed away entrance cryptographic contraptions for hard cross area issues, and constructed a secret entry limit of character based encryption plot. In the composing, Agrawal and Boyen introduced an IBE plan considering troublesome issues in networks and it was shown secure in standard model. Chen et al. proposed an IBE scheme from cross segments while the key can be repudiated capably.

Also, various HIBE plans have been proposed. Waters proposed one more security affirmation methodology for encryption plans.

The IBE and HIBE plans are under decisional Bilinear Diffie-Hellman (DBDH) and decisional Linear (DL) doubts. In the composing, Lewko and Waters arranged a totally safeguarded HIBE contrive, which can make a short ciphertext, and the arrangement is turned out to be secure under three static speculations.

Though some person based encryption plans have been used on private prosperity records, they are not helpful on this current reality reduced contraptions. Ibraimi et al. proposed a character set up middle person re-encryption plot regarding clinical consideration. Their ideal relies upon BF-IBE's thought. In their arrangement, a ciphertext under someone's public key can be re-encoded on a middle person server, and server can yield a new ciphertext under other client's public key. Wang et al. proposed a patient-driven cloud-based secure PHR system. In their arrangement, the patient's PHR can be taken care of in the cloud organization securely, and the serious characters have the right agree to disentangle the encoded PHR data. Following the possibility of different evened out character based encryption, Zhang et al. proposed a security protecting sharing of PHR in the cloud. In any case, as of not long ago, these plans need significant computation, for instance, bilinear aide estimation, which are

not sensible for the greater part helpful devices in private prosperity records sharing systems. Subsequently, we going to focus on the most capable technique to lessen the estimation cost for PHRs sharing system. Likewise, we are arranging a conveyed unscrambling plan in PHRs, which allows various get-togethers in a comparative office to interpret the ciphertext without reproducing the looking at private key.

III. PROPOSED WORK

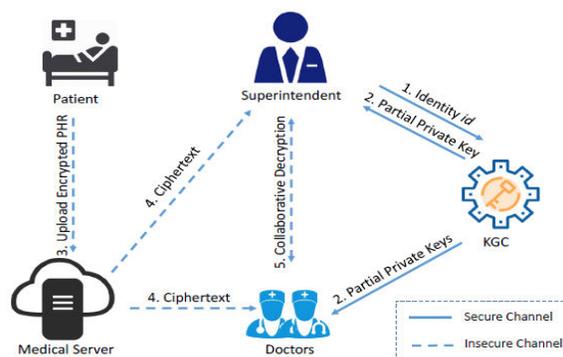


Fig. 3: System Model

In the above Figure, the electronic personal records sharing system model works as follows:

- 1) The superintendent sends the department's identity id to the KGC.
- 2) The KGC extracts and returns keys the partial private keys to the superintendent and the doctors of the department id.
- 3) The patient shares the PHRs with the department by encrypt the PHRs data under the department's identity id, then

upload the ciphertext to the medical server.

- 4) The superintendent and the doctors download the ciphertext from the medical server.
- 5) The superintendent and the doctors compute some temporary values by using their partial private keys.
- 6) After interacting with the doctors, the superintendent decrypts the ciphertext and outputs the PHRs.

Implementation Modules

- *Cloud server*

In this module, the cloud server offers computation & storage services of its users. The cloud server admin login to system and performs various operations & like view patients & authorize, view doctors & authorize them, view uploaded medical server files, view attackers, view transactions & view results

- *KGC*

In this module, KGC login to the system and distribute the keys to doctors.

- *Superintendent*

In this module, superintendent login to system & view the users ciphertext data.

- *Patient*

In this module, the patient register to the system and wait for the authentication after

authorization he login to system and perform various operations like view profile, upload patient health records before uploading encrypt the data then upload to cloud server, view any reports, view search controls decryption key permitted reports, verify report and delete report.

- *Doctor*

In this module, doctor register to system and wait for authentication, after authorized he login to system and perform various operations like view my profile, send such control request to KGC if KGC accepts his request he search the report & view the encrypted reports, if he wants the decrypt & download the reports he send decryption key request to KGC if KGC permitted then doctor view key response & download the patient health reports.

Implementation Algorithms

We present a novel distributed decryption scheme based on BF-IBE for electronic PHRs sharing system. Our goal is to secure share the PHRs to multiple parties and prevent the doctor's private key from being compromised. Firstly, the KGC invokes Setup to get the public parameters params and master secret key. Then, the doctor or the department registers to the system to obtain the partial private keys.

User Register

- 1) The department sends its identity id to KGC.
- 2) When KGC received the identity id, it extracts the partial private keys for the superintendent and the doctors by using the master secret key and params.
- 3) KGC sends $D1$ to the superintendent, and s_i ($1 < i < n$) to the doctors. The KGC extracts and distributes the partial private keys to the superintendent and the doctors.

Encrypt and Upload

- 1) Determines the identity of the shared department.
- 2) Invokes the original Encrypt algorithm of BF-IBE to encrypt the PHRs M to get the ciphertext $C = (C1;C2;C3)$.
- 3) Uploads his/her encrypted PHRs C to the medical server.

Download and Decrypt

When the doctors of the department need to access the user’s PHRs, they should interact with the superintendent. Finally, the superintendent outputs the PHRs.

- 1) The superintendent and the doctors download the encrypted PHRs $C = (C1;C2;C3)$ from the medical server.
- 2) Compute partial decryption values
- 3) Generate the zero-knowledge proofs and send to superintendent

- 4) The superintendent receives partial decryption values and proofs verifies.
- 5) Decrypt the ciphertext
- 6) Check the validation of the ciphertext
- 7) Output and share the PHRs

IV. RESULTS

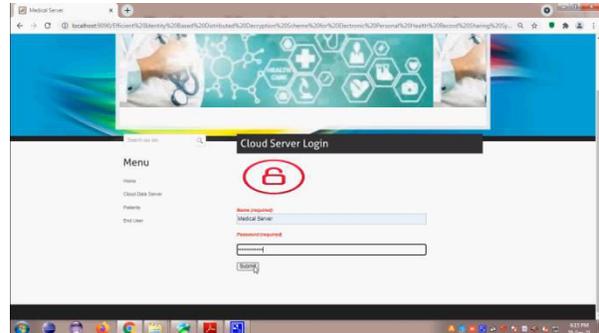


Fig. 4: Cloud Server Login



Fig. 5: View Medical Server Files

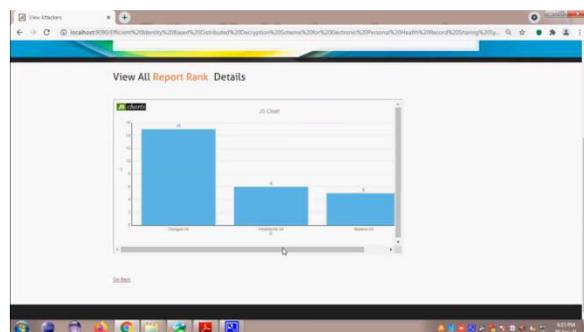


Fig.6: View All Reports Rank

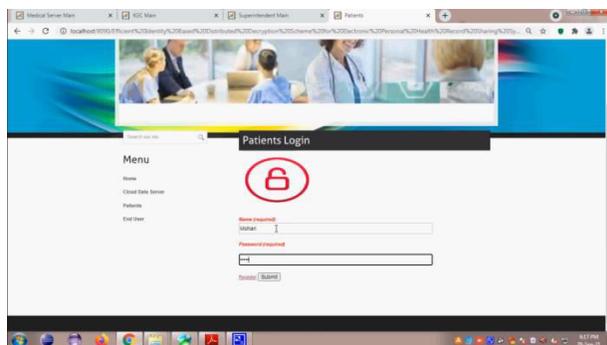


Fig. 7: Patient Login

V. CONCLUSION

Electronic Personal Health Record Sharing Systems are being widely used. Security and privacy issues are becoming critical in such systems and environments. Securing sensitive data of users, such as medications, chronic health problems, immunization history and the private keys in these environments is crucial and challenging. We employed the Boneh Franklin identity-based encryption scheme to design an efficient and secure e-health personal health record sharing system in this paper. In our proposed scheme, patient can encrypt the PHRs under the identity of a doctor or a department. The ciphertext can be decrypted securely by multiple parties. Specifically, our scheme is lightweight for the mobile devices, and it allows the parties to decrypt the ciphertext without reconstructing the private key.

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